

The importance of algorithms in Computer Science

- or -

how Manao learned how to defeat all the monsters.

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Computers are everywhere





Computer Science



Algorithms



Algorithms = Human (high-level) way of teaching a computer what to do to solve a problem



One algorithm can change the world







One algorithm can change the world



Why was the Google algorithm successful?

Solved a problem – search – in an *elegant*, *efficient*, *correct* and *scalable* way.

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Manao and the monsters



- Manao is traversing a valley inhabited by monsters.
- During his journey, he will encounter several monsters.
- The scariness of each monster is a number dread/scare.
- Manao is not going to fight the monsters (watje).



Manao and the monsters



- Instead, he will bribe some of them and make them join him.
- The monsters are not too greedy (price = 1 or 2).



Manao and the monsters



- Manao meets monster: bribe or not bribe?
- No bribe, not scary enough: monster will attack!
- Bribe: monster will join the party!
- So: if monster is stronger, Manao bribes.
- If monster is not strong: choice.



Manao and the monsters: example



- Dragon: bribe (always for the first monster).
- Hydra: Manao can either pass or bribe her.
- Killer Rabbit:
 - ① Bribed Hydra: total scariness = 13 > 10. Pass the rabbit! Total cost = 1 + 1 = 2 coins.
 - **2** Bribe the Rabbit. Total cost = 1 + 2 = 3 coins.



- There are between 1 and 20 monsters.
- Their scariness level is between 1 and 2,000,000,000, inclusive.
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Goal

Minimize the amount of money Manao needs to pay.

- Monster 0: no option, we bribe.
- Monster 1: two options.
 - **1** Not bribe: we are scary enough, money and scariness \iff .
 - $oldsymbol{2}$ bribe: scariness and money spent \uparrow .

:

 Monster p: we have passed the first p monsters and we have a total scariness level of currentScare.

How do we decide what to do?



- Monster p: Bribe or scare?
 We will compute the minimum price at step p given how scary we are M(p, CurrentScare) by analyzing our two options:
 - **1** bribe: it will cost us $price_p$, we earn $scare_p$.

$$P_1 = price_p + M(p + 1, CurrentScare + scare_p)$$

2 scare: $currentScare > scare_p$, not pay, not scarier.

$$P_2 = M(p+1, CurrentScare)$$

Choose best option: $M(p, CurrentScare) = min(P_1, P_2)$.

In the solution:

 We were using the minimum price of the next monster – recursive algorithm.

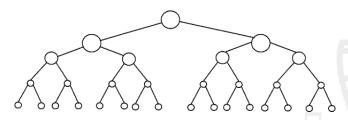
In the solution:

- We were using the minimum price of the next monster recursive algorithm.
- When do we stop: p = n all monsters done, price is 0.



How efficient is this algorithm?

Let's measure how many options we have:



At each step we double our options. For n monsters we have

$$2 \times 2 \times \cdots \times 2$$

n times

How efficient is this algorithm?

$$\underbrace{2 \times 2 \times \cdots \times 2}_{n \text{ times}} = 2^n$$

For 20 monsters we have $2^{20} \sim 1.000.000$.



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Good or bad?















B, dread = 2^10



E, dread = 10

F, dread = 10





M, dread = 10









R, dread = 10

L, dread = 10





T, dread = 10

P, dread = 10

There can be now up to **50** monsters

P. dread = 10



Is our previous algorithm ok?





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 $2^{50} = VERY BIG NUMBER$





Is our previous algorithm ok?

 $2^{50} = VERY BIG NUMBER$

Actually, with 2^{50} grains of rice you can cover the whole of the Netherlands with ca 2 meters of rice (calculations in Henk Barendregt's kitchen!).

Can we do better?





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Different perspective: what is the maximum price to pay for *n* monsters?



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Different perspective: what is the maximum price to pay for n monsters? $2 \times n$



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Instead of minimizing, ask...

For each price p from 0 to $2 \times n$, can we cross without paying more than p?

For $n \le 50$ there are at most 101 values for p and we can answer the problem *fast* (in polynomial time).



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solution = chocolate bar!

Topcoder



Manao and the monsters appeared in a programming match of TopCoder. Both problems were solved in under 4 minutes!



Lessons to take home

- The parameters of the problem matter.
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A lot of Computer Science is like solving puzzles!

